

# Introduction to Security

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- Problems
- General Security
- Cryptography & Protocol reviewed

- Life was beautiful before computer, getting worse after Internet
- Distributed systems: each depends on others
- How can we assure system behaves correctly & securely ?
- Can we trust systems on the other side ?

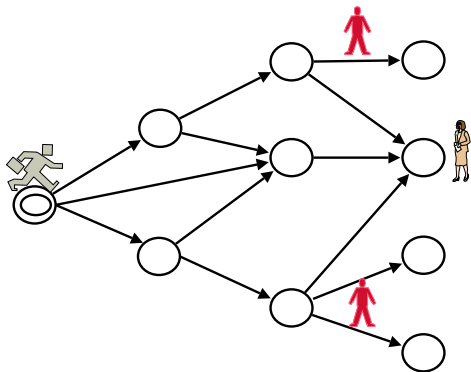


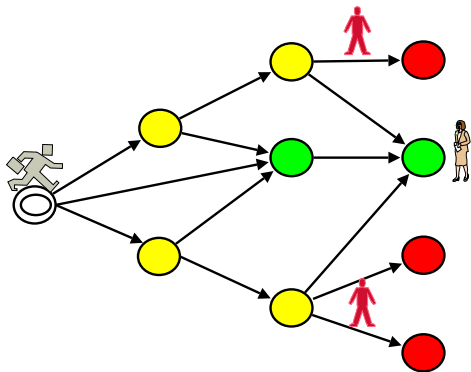
- How to assure security of the network
- How to quantify risk
- What are the boundaries of the system ?
- Relationship to political, social, enomical aspects is not well understood
- Uncertainty

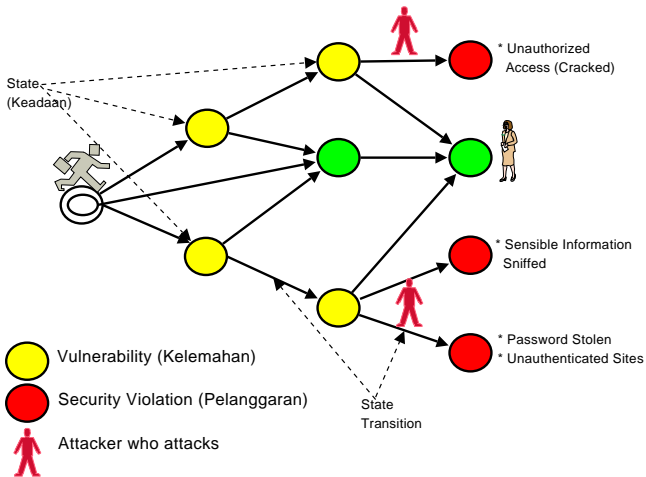


- In General :
  - Computer Security deals with the method against unauthorized actions in computer systems
- More General :
  - Dependability to other institution is also insecurity











# What is secure ?

- Computer based system behaves according to
  - 1 algorithms (program/software)
  - 2 user direction (input)
- Given systems & environment:
  - 1 **secure states** (system maintains security properties)
  - 2 **insecure states** (violation of security policy)
  - 3 **paths lead to insecure states** (vulnerability)
- Correctness: maintain intended behavior according to correct specification while unintended behavior will not be reachable
- Security policy: definition of (1,2,3)
- Attack : **Persistence**, **Intentional**, Outsider vs Insider, Loss vs Gain



- Towards vulnerable system [Abadi] :
  - Interaction with uncertain physical, network, software environment.
  - Using public network, distributed administration, diverse operators
  - COTS, business demand is the priority, Monocultures



- Attack :
  - Physical attack  
ex. theft of hddisk/cdroms, bombing, etc.
  - Syntatic attack  
ex. buffer overflow, domain theft, SQL injection
  - Semantic attack  
ex. Social engineering, site phising



- We define secure states and insecure states
- Define paths which always bring system in secure states: what is permitted
- Define also paths might lead to insecure states: what is not allowed
- Specified in formal language for clearness, unambiguity, consistency, and verifiability
- Written in natural language for better understanding



- Confidentiality (Secrecy)  
Unauthorized disclosure of information is not reachable  
(Access Control - Cryptography)
- Integrity  
Unauthorized modification of information is not reachable
  - Data integrity - Origin integrity / authentication
  - Prevention (access control) - Detection (hash function)



- Availability  
Prevention & detection of denial of service
- Accountability  
The availability and completeness of the track of past system states  
Who - Whom - What - When - Where  
Implementation: should be forensic ready



- Prevention  
to fail the attack
- Detection  
to detect unprevented attack
- Recovery  
to stop the attack & repair attacked system

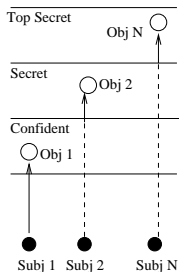
- Way to enforce security policies
- How to limit system behavior according to policies
- Specification - Design - Implementation - Operation & Maintenance - Audit (Forensic)
- Access Control & Cryptography



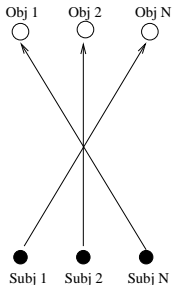
- Reference Monitor
- Set of precise rules according to security policy, applied as a filter to the transition states of the system, which prevents system in entering insecure state
- Authentication is mandatory
- **Subjects, Objects, Actions**, Time, Location, and other attributes

# Model of Access Control

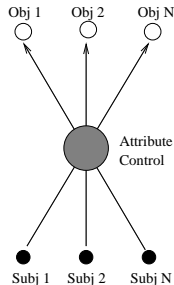
**Mandatory Access Control (MAC)**



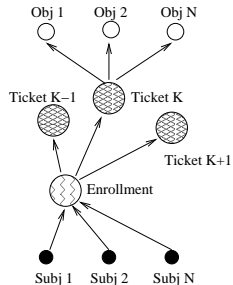
**Discretionary Access Control (DAC)**



**Role Based Access Control (RBAC)**



**Safe Dealing (SD)**



# Implementation: Cryptography

- Algorithm to protect *secrecy* of data
- Also used to gain :
  - *authentication*
  - *integrity*
  - *non repudiation*
- Includes : *algorithm* and *key(s)*



- $Chipertext = Encrypt(Message, Key)$
- $Message = Decrypt(Chipertext, Key)$
- $Decrypt(Chipertext)$  hard without  $Key$
- Research questions :  
Is there any algorithm which is hard to compute original message but easy to verify it
- In implementation requires a protocol (*Cryptographic Protocol*)



$$A \mapsto B : \{M\}_{K_{ab}}$$

**Principal  $A$  sends  $B$  message  $M$   
encrypted with *shared-key*  $K_{ab}$**

- Key is shared between 2 *principals*
- Needs  $N^2$  keys for  $N$  principals
- Fast but *key management* is not easy
- Example of Cipher: DES, 3DES, Blowfish, AES



$$A \mapsto B : \{M\}_{\mathcal{K}_b}$$

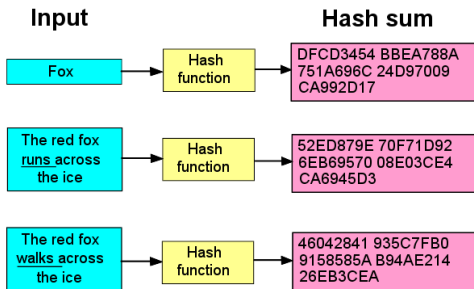
**Principal  $A$  sends  $B$  message  $M$   
encrypted with  $B$ 's public-key  $\mathcal{K}_b$**

- Only with *private-key*  $\mathcal{K}_b^{-1}$ ,  $B$  can *decrypt*  $M$
- Principal has its own  $\mathcal{K}$  which is *published* and  $\mathcal{K}^{-1}$  which must be kept *secret*
- Key management is less difficult, usually managed by *Certification Authority*
- Example of Cipher: RSA (Rivest-Shamir-Adleman), Elliptic-Curve



# One-Way-Hash

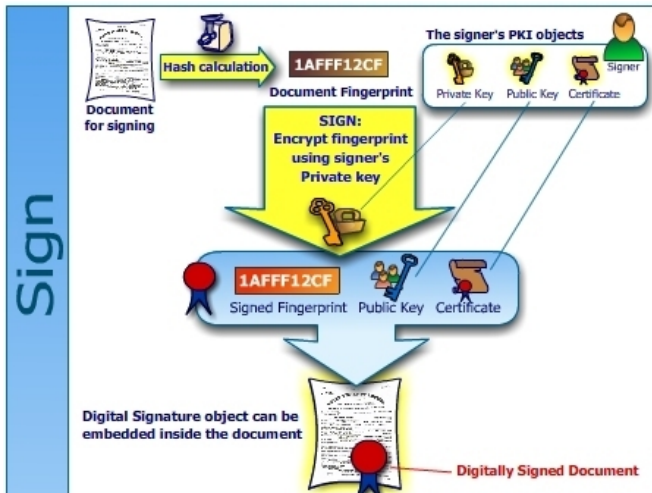
- Algorithm to compute large data into small integer, producing *fingerprint of the message*



- Used for maintaining **integrity** of message being transferred
- Example: MD5, SHA1, SHA-256, Ripemd, Haval

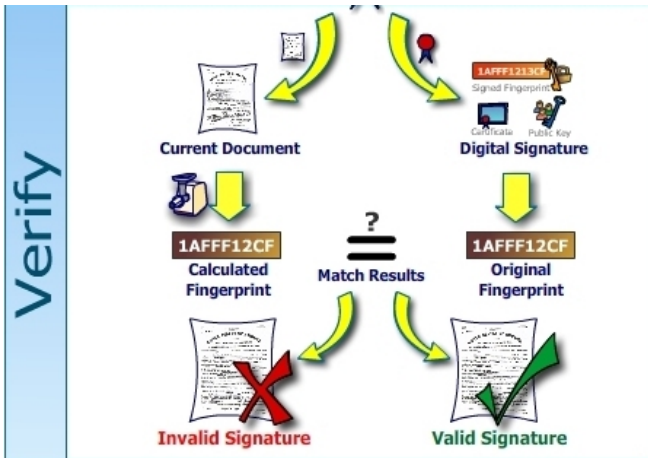


# Digital Signature (Sign)





# Digital Signature (Verify)

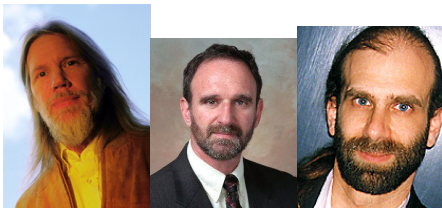


$$A \mapsto B : \{M, \{Hash(M)\}_{K_a}\}_{K_b}$$

- A's digital signature on a message is the hash of message encrypted with A's private-key
- Authentication: only with A's public-key, the hash can be decrypted
- Integrity: Hash function
- Confidentiality: message can be decrypted only with B's private-key
- Non-Repudiation: explain for your self



# Pictures of Cryptographer



- Implementation of Cryptography Algorithm
- Achieving *security properties* (authentication, secrecy, etc.)
- Example :
  - **Needham-Schroeder** (*authentication*)
  - **Kerberos** (*authentication*)
  - **SSL/TLS** (*auth - secrecy*)



- Example : **Needham-Schroeder Protocol**

**M1**  $\mathcal{A} \mapsto \mathcal{S} : A, B, N_a$

**M2**  $\mathcal{S} \mapsto \mathcal{A} : \{N_a, B, K_{ab}, \{K_{ab}, A\}K_{bs}\}K_{as}$

**M3**  $\mathcal{A} \mapsto \mathcal{B} : \{K_{ab}, A\}K_{bs}$

**M4**  $\mathcal{B} \mapsto \mathcal{A} : \{N_b\}K_{ab}$

**M5**  $\mathcal{A} \mapsto \mathcal{B} : \{N_b - 1\}K_{ab}$

- Introducing Nonce ( $N$ )



- More example : **Kerberos Protocol**

**M1**  $\mathcal{A} \mapsto \mathcal{S} : A, B$

**M2**  $\mathcal{S} \mapsto \mathcal{A} : \{T_s, L, B, K_{ab}, \{T_s, L, K_{ab}, A\}K_{bs}\}K_{as}$

**M3**  $\mathcal{A} \mapsto \mathcal{B} : \{T_s, L, K_{ab}, A\}K_{bs}, \{A, T_a\}K_{ab}$

**M4**  $\mathcal{B} \mapsto \mathcal{A} : \{T_a + 1\}K_{ab}$

- Introducing TimeStamp ( $T$ ) and Lifetime ( $L$ )
- Used in many system, including Windows



## Problem :

- Wrong design could lead to flaw
  - Needham-Schroeder Protocol
  - SSLv1.0
- Wrong implementation could lead to vulnerability
  - Padding problem in SSL, SSH, and WTLS
  - User Interface design in Browser
- Vulnerability arise between two protection technologies (Anderson, Ross)



# Assurance : Formal Method

- To prove correctness in achieving security properties which protocol carry out
- There are two development approach :
  - Extention from method used in communication
  - Newly developed method
- Four classifications :
  1. General purpose tools
  2. Logic based
  3. Expert System
  4. Algebraic approach

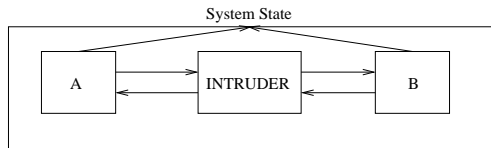




# Formal Method

## Using General Purpose Tools

- Treated as ordinary comm. protocol
- Adversary is explicit, capable in read, intercept, and modify messages
- Method : FSM, CSP, FDR, Petri Nets  
Example : Lotos, Ina Jo, Murphy



- Investigate every possible scenario of ***Attack - Flaw - Defence***
- Needs to define **insecure states** and *search* paths to them
- More successful than General Purpose Tools
- Example : Interrogator by Millen, NRL Protocol Analyzer by Meadows, Longley and Rigby



- Capabilities in **modeling knowledge** which represents component in cryptographic operation (*Nonce, Key(s), and old messages*)
- Example :
  - Dolev - Yao (term re-writing systems)
  - Sphi - Calculus by Abadi and Gordon (to prove secrecy)



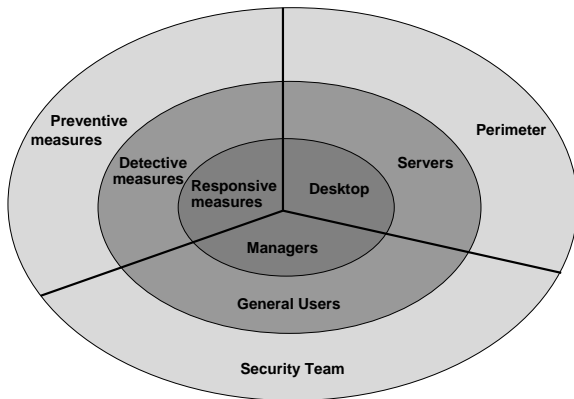
- One sees crypt. protocol as **distributed algorithm**
- Develop logics from **modal logic**
- There are **inference rules**
- Goal is to derived statements which **represents correct condition**
- Example : *BAN Logic and GNY Logic*



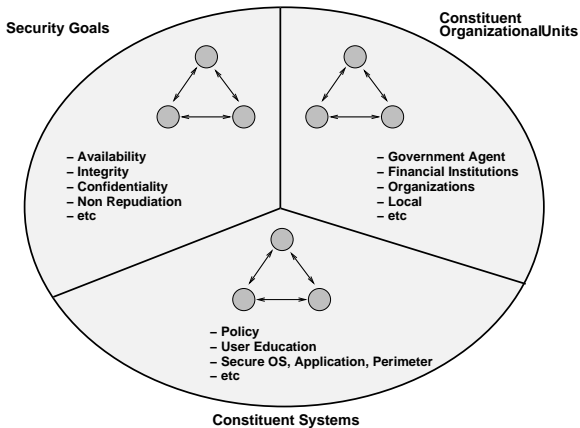
# Towards Secure System

- Specification : Security Policy
- Implementation : Security Mechanism
- Correctness : Assurance
- Man - Machine - Management

# Towards ... (cont)



# Towards .. (cont)



# Books, Papers, and Links

- Ross Anderson, “Security Engineering”
- Matt Bishop, “Computer Security”
- Schneider et. al. “Modelling and Analysis of Security Protocols”
- Martin Abadi’s homepage at *<http://www.cse.ucsc.edu/~abadi>*





End of this presentation